



Setting the Stage

The Idea of Epistemic Logic

Epistemic logic is that branch of philosophical logic that seeks to formalize the logic of discourse about knowledge. Its object is to articulate and clarify the general principles of reasoning about claims to and attributions of knowledge—to elucidate their inferential implications and consequences. In pursuing this goal, it deals principally with propositional knowledge (along the lines of “Smith knows *that* coal is black”) and secondarily also with interrogative knowledge (along the lines of “Jones knows *where* the treasure is buried and *who* put it there”).¹ It is the object of this book to give an overview of the discipline by setting out in a formalized manner the general principles for reasoning about such matters.

The History of Epistemic Logic

Epistemic logic is a product of the second half of the twentieth century. After the preliminary work of Rudolf Carnap’s deliberations about belief sentences,² epistemic logic was launched in an important 1948 paper by the Polish logician Jerzy Łoś.³ Łoś developed what he called a logic of “belief” or “acceptance” based on an operator Lxp for

“the individual x believes (or is committed to) the proposition p ,” for which he stipulated axiomatic rules substantially akin to those to be specified for the knowledge operator to be introduced below. A spate of publication during the 1950s by such logicians as Alonzo Church, Arthur Prior, Hilary Putnam, G. H. von Wright, and the present writer further extended the range of relevant deliberations. During the 1960s various authors carried matters forward, with the first book on the topic (by Jaakko Hintikka) appearing in 1962. Since then there has been a small but steady stream of work in the field. (For details, see the bibliography.)

Fundamentals of Notation

The present treatment of epistemic logic undertakes the construction of a deductively formalized system s that is adequate for this purpose. As is often the case with axiomatic treatments, the present discussion illustrates how a modest set of basic assumptions provides a great deal of instructive information about the conceptual anatomy of the idea at issue.

Use will be made here of the familiar resources of propositional and quantificational logic supplemented by the machinery of quantified modal logic. All of the following symbols will accordingly be used in the standard way:

- \sim , $\&$, \vee , \supset , and \equiv for the familiar propositional connectives
- \forall and \exists for universal and existential quantification
- \square and \diamond for the modalities of necessity and possibility

Additionally, the following notational conventions will be employed:

- Kxp (“ x knows that p ”)
- K^*xp (“ p is derivable from propositions that x knows”)
- x, y, z, \dots as variables for knowers: intelligent individual (or possibly groups thereof)
- p, q, r, \dots as variables for propositions (contentions to the effect “that such-and-such is the case”)
- t, t', t'', \dots as variables for specifically true propositions (note that

$(\forall t)F(t)$ amounts to $(\forall p)(p \supset F(p))$ and $(\exists t)F(t)$ amounts to $(\exists p)(p \& F(p))$, F being an arbitrary propositional function
 u, u', u'', \dots as variables for objects of consideration or discussion
 F, G, H, \dots as variables for properties or features objects or of propositions
 S, S', S'', \dots as variables for sets of objects or propositions
 Q, Q', Q'', \dots as variables for questions

The quantificational logic at work here is type differentiated: p, q, r , and so on stand for propositions x, y, z , and so on for knowers and so on. One could, in theory, employ a single class of variable α, β, γ and so on, and then render

$(\forall x)Fx$ as $(\forall \alpha)(\alpha \in K \supset F\alpha)$,
 where K represents the set of knowers,
 $(\exists p)Gp$ as $(\exists \alpha)(\alpha \in P \& G\alpha)$,
 where P represents the set of propositions,

and the like. But this more elaborate style of presentation would make our formulations needlessly complicated and less easily read. It should be noted that the various domains at issue (knowers, propositions, truths, and the like) are all nonempty, so that the inference from “all” to “some” is appropriate in all cases.

Certain special symbols will be employed as follows:

$\vdash p$ for “ p is a thesis of our system (s)”

$\Vdash p$ for $\vdash (\forall x)Kxp$

$p \vdash q$ for $\vdash p \supset q$

$p \Vdash q$ for $\vdash (\forall x)Kxp \supset Kxq$

$p @ Q$ for “ p answers the question Q ”

The symbol \vdash will also be called upon to serve as an index of entailment through the following equivalence:

$p \vdash q$ iff $\vdash (p \supset q)$

Since the antecedent p may disaggregate into the conjunction of a series of propositions, p_1, p_2, \dots, p_n , this stipulation renders our system

subject to what is standardly called the deduction theorem, on the basis of the following equivalence:

$$p_1, p_2, \dots, p_n \vdash q \text{ if and only if } p_1, p_2, \dots, p_{n-1} \vdash p_n \supset q.$$

With propositional knowledge of matters of fact, the basic unit of assertion will be a statement of the form “ x knows that p ” (Kxp). Such propositional knowledge is a matter of a relationship—a *cognitive* relationship—between a person and a true proposition. And just as for an otherwise unidentified individual x one can uniformly substitute the name of any individual, so for an otherwise unidentified proposition p one can uniformly substitute any other. The use of variables thus affords a gateway to generality by providing for substitution. For example, since it obtains as a general principle that

If Kxp , then p ,

one automatically secures a vast range of such other assertions as

If $Ky(p \ \& \ q)$, then $p \ \& \ q$,

which results from the preceding via the substitutions y/x and $(p \ \& \ q)/p$.

Recourse to symbolic representation enables us to achieve greater precision. For instance, in ordinary language “ x does not know that p ” is equivocal as between $\sim Kxp$ and $p \ \& \ \sim Kxp$, which would be more accurately formulated as “ p , and x does not know it.”

Theses of the System

As already mentioned, \vdash here serves as an assertion symbol indicating that what follows qualifies as a general principle of the system of epistemic logic (s) that is under construction. By convention its employment conveys implicit universality for any free variables. Thus,

$$\vdash Kxp \supset p$$

asserts that $(\forall x)(\forall p)(Kxp \supset p)$ holds in our system. A proposition that qualifies as a thesis of the system should be seen as being true on logico-conceptual grounds alone. Its validation will rest entirely on

the specification of the terms of reference that are employed and thus on the conventions of meaning and usage that are being adopted. These theses accordingly serve to specify the conception of *knowledge* that is to be at issue. And since a “logic” of knowledge must deal in general principles, it is the establishment or refutation of such conceptually grounded generalizations that concern us at present. What is at work here is in fact a somewhat delicate reciprocal feedback process. A certain particular conception of knowledge guides the construction of our epistemic system. And the theses of this system define and precisify the particular conception of knowledge that is at issue.

In dealing with knowledge and its “logic” we are not, of course, functioning in a realm of total abstraction, as would be the case with “pure” (rather than applied) mathematical or theoretical logic. Instead, we are dealing with the resources of intelligent beings (not necessarily members of *Homo sapiens*) operating substantially within the limits imposed by the realities of this world of ours. Accordingly, the “facts of life” that reflect the cognitive situation of such beings and the conditions that define their situation in this world represent the ultimately factual (rather than purely theoretical) circumstances that a logic of knowledge as such will have to reflect. In particular, knowers have to be construed as finite beings with finite capacities, even though reality, nature, has an effectively infinite cognitive depth in point of detail, in that no matter how elaborate our characterizations of the real, there is always more to be said.⁴ The reality of it is that epistemic logic is an applied logic and its theses, being geared to salient feature of the established concept of knowledge, stand correlative to the ways in which we actually do talk and think about the matter.

Propositions as Objects of Knowledge

There is nothing problematic about saying “ p , but x does not know (or believe) it.” But in the special case of $x =$ oneself (the assertor), this otherwise viable locution is impracticable. This discrepant state of affairs has become known as “Moore’s paradox” after G. E. Moore, who first puzzled over it.⁵ Of course, there would be nothing amiss about saying “I surmise (conjecture, suspect) that p but do not actually know

(or confidently believe) it.” But in making a flat-out, unqualified statement we stand subject to the ground rule that this purports knowing the truth of the matter, so that in going on to add “but I do not know (or believe) it” to an assertion of ours, we take the inconsistent line of giving with one hand what we take away with the other. Our categorical (that is, unqualified) assertions stand subject to an implicit claim to truth and knowledge, and we thus authorize the inference from asserting p both to Kip and to p itself. Accordingly, when our system s is held to make an explicit assertion, this will be something that we ourselves purport to know, so that we then have it that $\vdash p$ entails $(\exists x)Kxp$.

In general, claims to knowledge regarding individual objects or collections thereof can be reformulated with the machinery of propositional knowledge by means of quantification. Thus, consider

“ x knows the identity of Jack the Ripper”:

$(\exists p)(p \text{ identifies who Jack the Ripper was} \ \& \ Kxp)$

“ x knows the major features of London’s topography”:

$(\forall p)(p \text{ states a major feature of London’s topography} \ \supset \ Kxp)$ ⁶

Such statements about someone’s knowledge of individual objects can be reduced to propositional knowledge by employing either

$(\exists p)(p @ Q \ \& \ Kxp)$

or

$(\forall p)(p @ Q \ \supset \ Kxp)$

when $p @ Q$ abbreviates “ p answers the question Q .”

By and large, propositional knowledge represents a resource by whose means the other principal versions of the concept of knowledge can be recast and represented. However, some knowledge is not propositionally reducible, specifically, know-how of a certain sort. For we have to distinguish between

performatory know-how: x knows how to do A in the sense that x can do A ; and

procedural know-how: x knows how A is done in the sense that x can spell out instructions for doing A .

The second sort of know-how is clearly a matter of propositional knowledge—that x knows that A can be done by doing such-and-such things; for example, x knows that people swim by moving their arms and legs in a certain cycle of rhythmic motions. But, of course, x can know how A is done without being able to do A —that is, without x having the performatory skills that enable x to do A . (For example, x may know *that* a certain result is produced when a text is translated from one language to another without actually knowing *how* to make such a translation.) And, therefore, while propositional reduction is practicable with respect to *procedural* know-how, such a reduction will not be practicable with respect to *performatory* know-how, seeing that people are clearly able to do all sorts of things (catch balls, remember faces) without being able to spell out a process or procedure for doing so.⁷

All the same, the different modes of knowledge are inextricably interconnected. To know (propositionally) *that* a cat is on the mat one must know (adverbially) *what* a cat is. And this knowledge rests on knowing how to tell cats from kangaroos.